

SHIP: Scalable Hierarchical Power Control for Large-Scale Data Centers

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Introduction

- Data centers are expanding to meet new business requirement.
 - **Cost-prohibitive** to expand the power facility.
 - Upgrades of power/cooling systems lag far behind.
 - Example: NSA data center
- Power overload may cause system failures.
 - Power provisioning **CANNOT** guarantee exempt of overload.
 - Over-provisioning may cause unnecessary expenses.



Power control for an entire data center is very necessary.

Challenges

- **Scalability**: One centralized controller for thousands of servers?
- **Coordination**: if multiple controllers designed, how do they interact with each other?
- **Stability** and **accuracy**: workload is time-varying and unpredictable.
- **Performance**: how to allocate power budgets among different servers, racks, etc.?

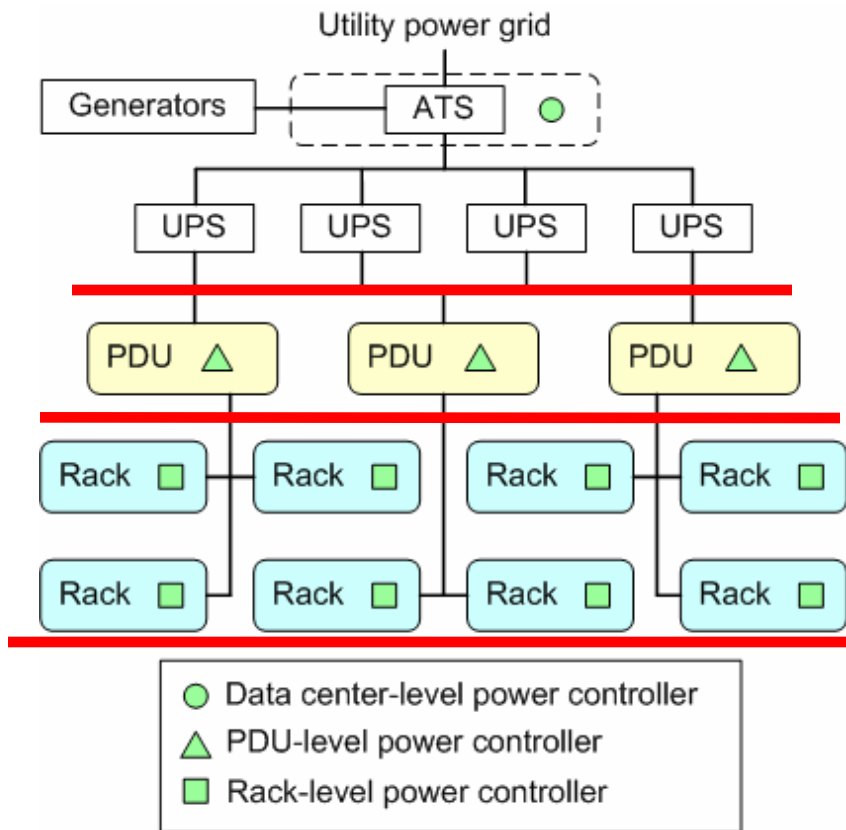
State of The Art

- Reduce power by improving energy-efficiency : [Lefurgy], [Nathuji], [Zeng], [Lu], [Brooks]
 - Based on heuristic and NOT enforce power budget.
- Power control for a server [Lefurgy], [Skadron], [Minerick], a rack, [Wang], [Ranganathan], [Femal]
 - Cannot be directly applied for data centers.
- *No “Power” Struggles* presents a multi-level power manager. [Raghavendra]
 - NOT designed based on power supply hierarchy
 - NO rigorous overall stability analysis
 - Only simulation results for 180 servers
- Use power as a knob to control performance requirements in OS level. [Horvath], [Chen], [Sharma]

What is This Paper About?

- SHIP: a highly Scalable Hierarchical Power control architecture for large-scale data centers
 - **Scalability**: decompose the power control for a data center into three levels.
 - **Coordination**: hierarchy is based on power distribution system in data centers.
 - **Stability** and **accuracy**: theoretically guaranteed by Model Predictive Control (MPC) theory.
 - **Performance**: differentiate power budget based on performance demands, *i.e.* utilization.

Power Distribution Hierarchy

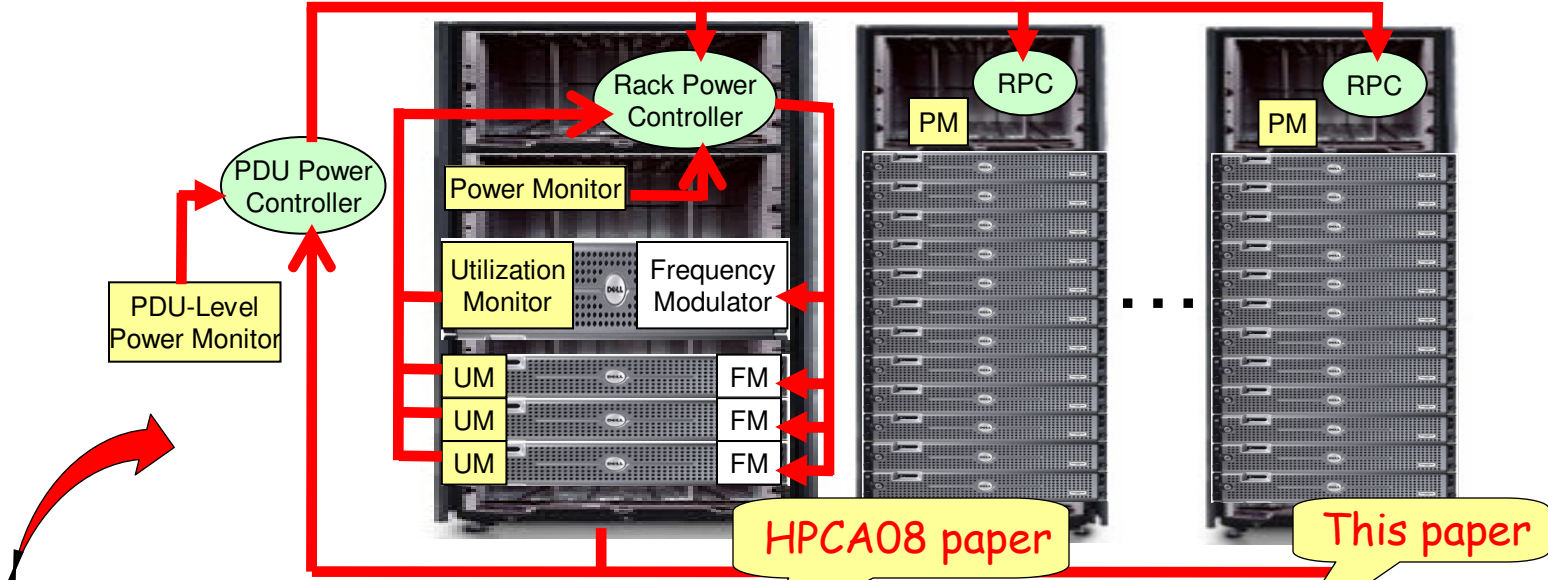


- A simplified example for a three-level data center

- Data center-level
- PDU-level
- Rack-level

- Thousands of servers in total

Control Architecture



| | Rack-level | PDU-level | Data center-level |
|----------------------|----------------------------------|-------------------------------|------------------------------------|
| Controlled variable | The total power of the rack | The total power of the PDU | The total power of the data center |
| Manipulated variable | The CPU frequency of each server | The power budget of each rack | The power budget of each PDU |

PDU-level Power Model

- System model:

$$pp(k+1) = pp(k) + \sum_{i=1}^N \Delta pr_i(k)$$

$pp(k)$: the total power of PDU
 $\Delta pr_i(k)$: the power change of rack i

- Uncertainties:

$\Delta pr_i(k) = g_i \Delta br_i(k)$ $\Delta br_i(k)$: the change of power budget for rack i
 g_i is the **power change ratio**.

- Actual model:

$$pp(k+1) = pp(k) + [g_1 \quad \dots \quad g_N] \begin{bmatrix} \Delta br_1(k) \\ \dots \\ \Delta br_N(k) \end{bmatrix}$$

Model Predictive Control (MPC)

- Control objective:

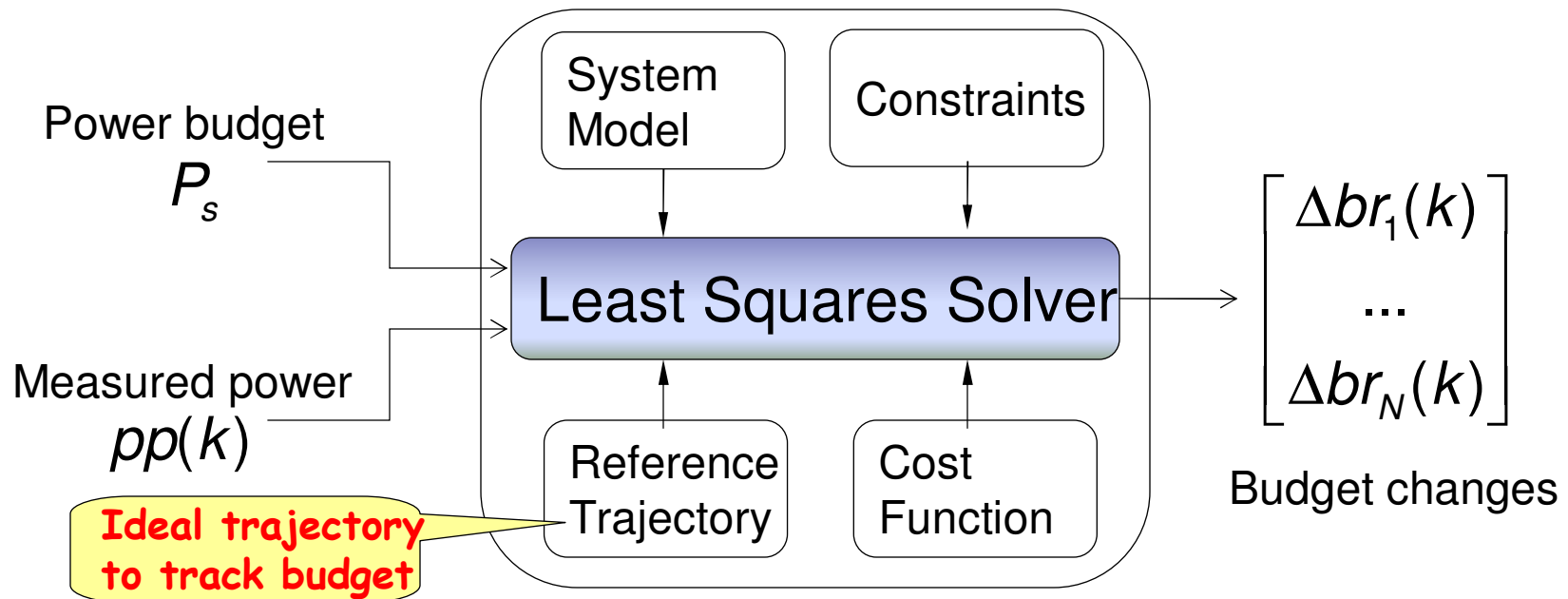
$$\min_{\{\Delta br_j(k) | 1 \leq j \leq N\}} (pp(k+1) - P_s)^2$$

$$\text{subject to : } P_{\min,j} \leq \Delta br_j(k) + br_j(k) \leq P_{\max,j} \quad (1 \leq j \leq N)$$

$$pp(k+1) \leq P_s$$

- Design steps:
 - Design a dynamic model for the controlled system.
 - Design the controller.
 - Analyze the stability and accuracy.

MPC Controller Design



$$V(k) = \sum_{i=1}^P \|pp(k+i|k) - ref(k+i|k)\|_{Q(i)}^2 + \sum_{i=0}^{M-1} \|\Delta br(k+i|k) + br(k+i|k) - P_{\max}\|_{R(i)}^2$$

Tracking error

Control penalty

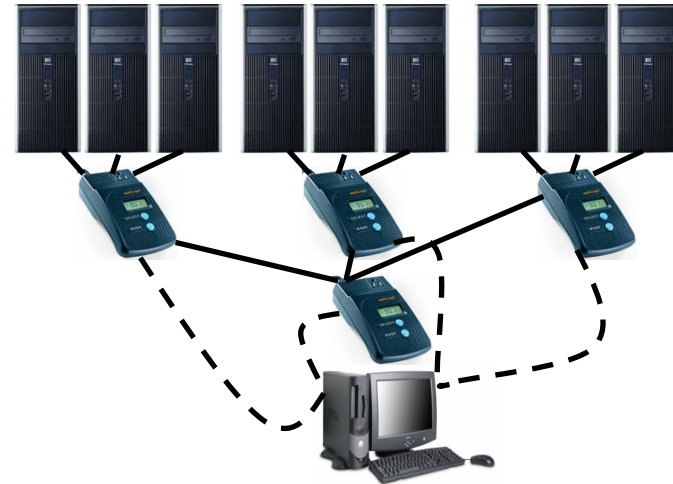
Stability

- Local Stability
 - g_i is assumed to be 1 at design time.
 - g_i is unknown a priori.
 - $0 < g_i < 14.8$: 14.8 times of the allocated budget
- Global Stability
 - **Decouple** controllers at different levels by running them in different time scales.
 - The **period** of upper-level control loop $>$ the **settling time** of the lower-level
 - Sufficient but not necessary

System Implementation

- Physical testbed

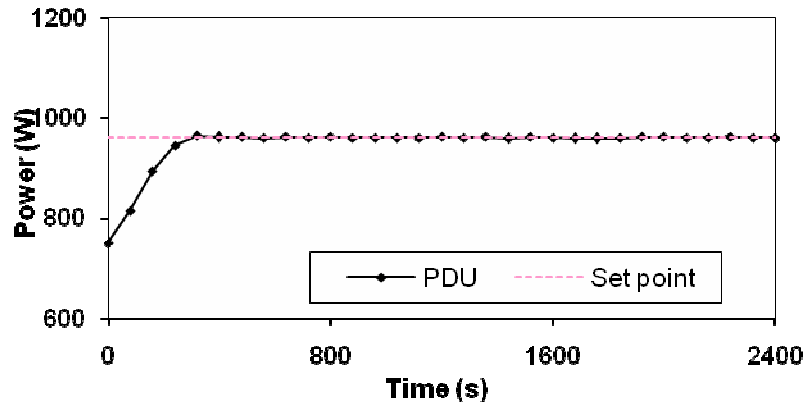
- 10 Linux servers
- Power meter (Wattsup)
 - error: $\pm 1.5\%$
 - sampling period: 1 sec
- Workload: HPL, SPEC
- Controllers:
 - call matlab function.
 - period: 5s for rack, 30s for PDU



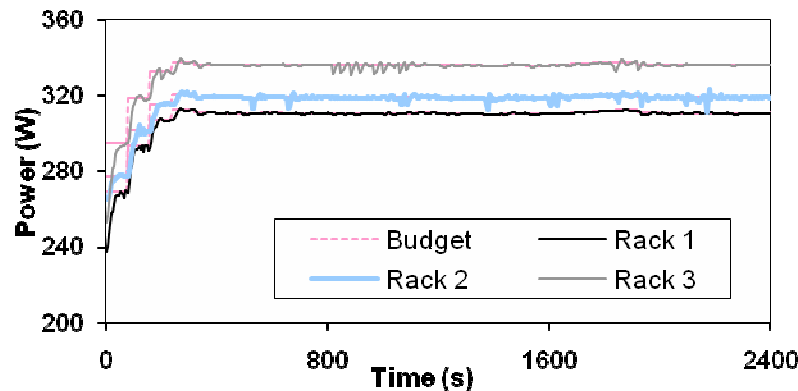
- Simulator (C++)

- Simulate large-scale data centers in three levels.
- Utilization trace file from 5,415 servers in real data centers
- Power model is based on experiments in servers.
- Generate 3 data center configurations.

Precise Power Control (Testbed)



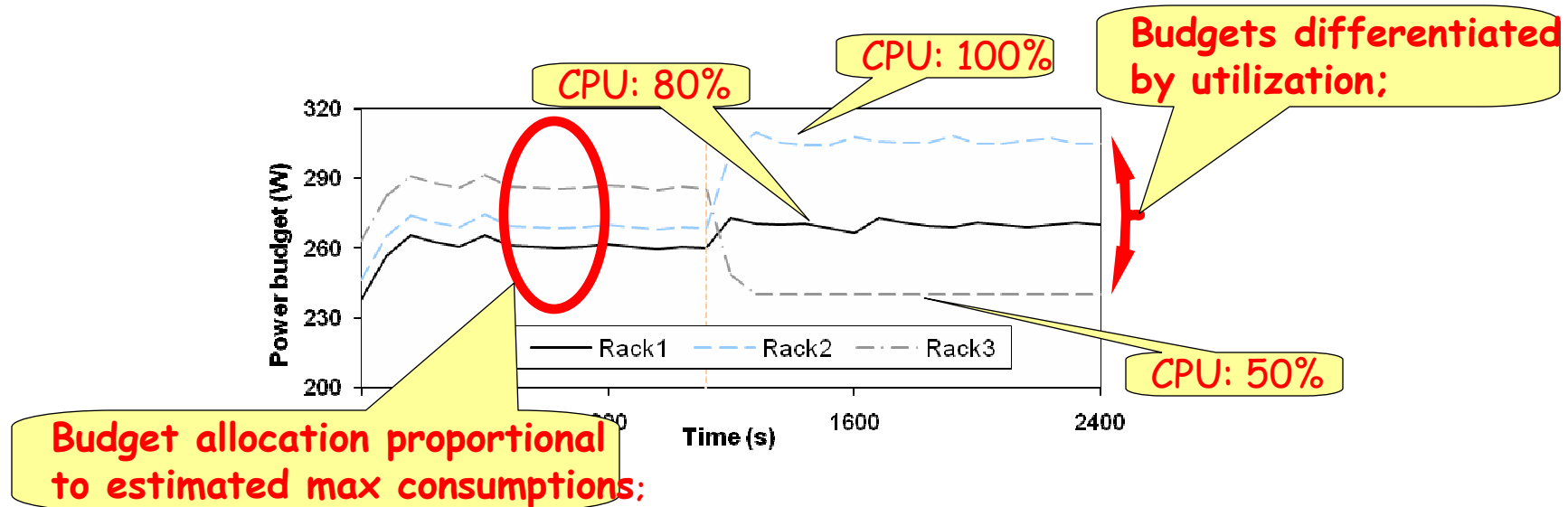
- Power can be precisely controlled at the budget.
- The budget can be reached within 4 control periods.



- The power of each rack is controlled at their budgets.
- Budgets are proportional to P_{\max} .

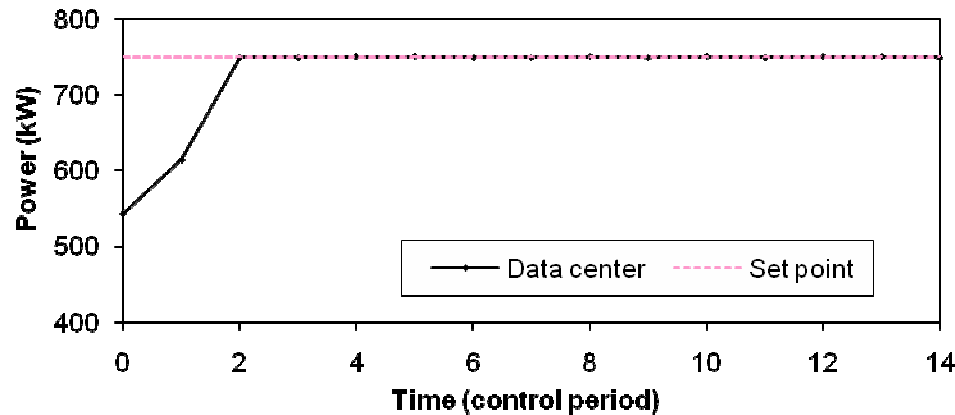
- Tested under other set points

Power Differentiation (Testbed)

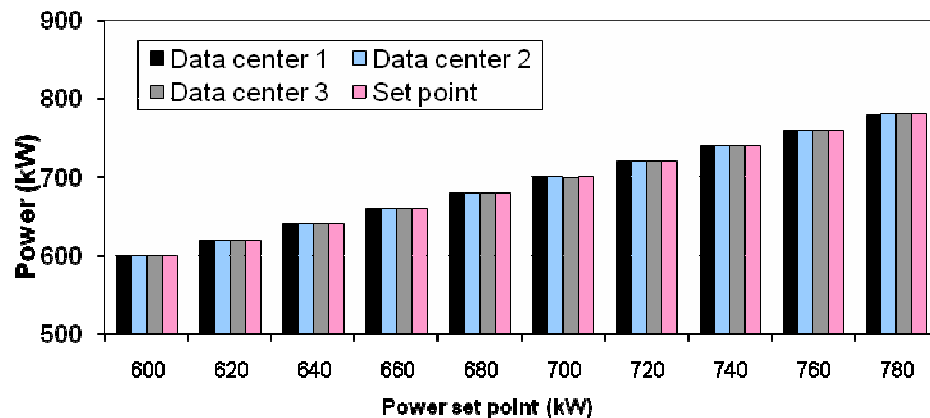


- Capability to differentiate budgets based on workload to improve performance
- Take the utilization as the optimization weights.
- Other differentiation metrics: response time, throughput

Simulation for Large-scale Data Centers

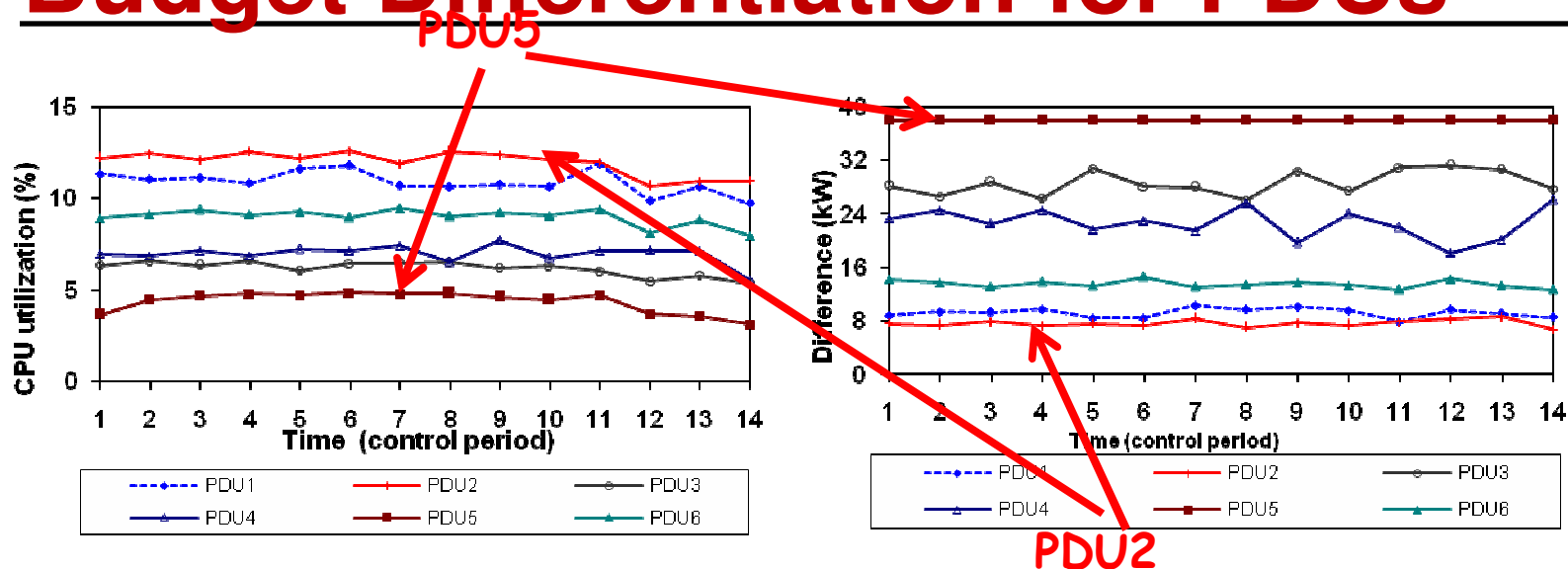


- 6 PDU, 270 racks
- Real data traces
- 750 kW



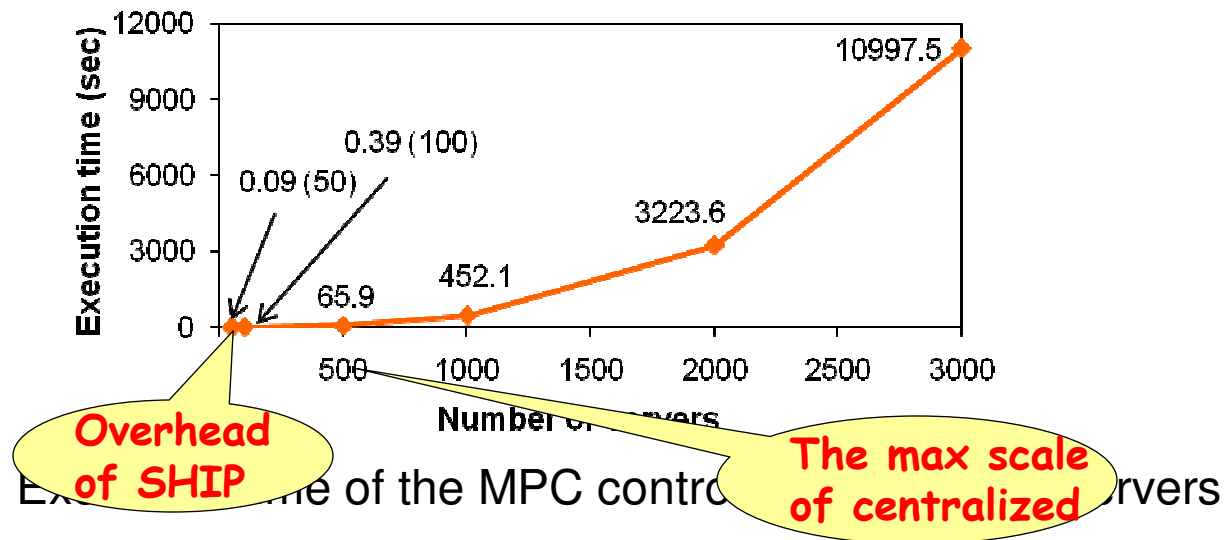
- Randomly generate 3 data centers
- Real data traces

Budget Differentiation for PDUs



- Power differentiation in large-scale data centers;
 - Minimize the difference with estimated max power consumption.
 - Utilization is the weight.
 - The **difference order** is consistent with the **utilization order**.

Scalability of SHIP



| | Centralized | SHIP |
|------------------------|-------------|----------|
| Level | One level | Multiple |
| Computation overhead | Large | Small |
| Communication overhead | Long | Short |
| Scalability | NO | YES |

Conclusion

- SHIP: a highly Scalable Hierarchical Power control architecture for large-scale data centers
 - **Three-levels**: rack, PDU, and data center
 - MIMO controllers based on optimal control theory (**MPC**)
 - Theoretically guaranteed **stability** and **accuracy**
 - Discussion on **coordination** among controllers
- Experiments on a physical testbed and a simulator
 - Precise power control
 - Budget differentiation
 - Scalable for large-scale data centers

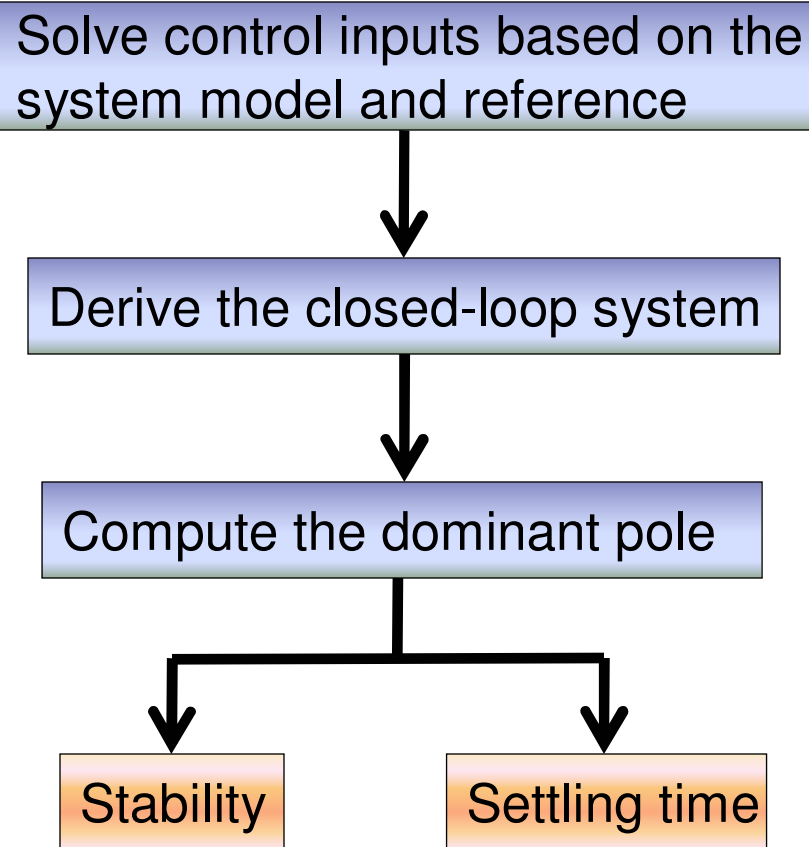
Acknowledgment

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Thank you!

Backup Slides

Stability Analysis



More Implementation Details

- CPU modulator
 - 4-5 frequency levels to scale
 - fraction levels:
 - For 2.8, that is: 2, 3, 3, 3, 3 with 5 subintervals.
 - 50 subintervals in each period of rack controllers
- Trace file
 - From 5415 servers in multiple data centers (manufacturing, financial, telecommunication, retail sectors)
 - Average CPU utilization every 15 minutes
 - From 00:00 on July 14 to 23:45 on July 20 in 2008

Reference Trajectory

$$ref(k + i | k) = P_s - e^{-\frac{T_p}{T_{ref}} i} (P_s - pp(k)), 1 \leq i \leq P$$

- T_p and T_{ref} specify the speed of system response.
- P : prediction horizon